# NEXT GENERATION INTERNET The GNU Taler Payment System

Christian Grothoff

December 2025

# Agenda

Motivation & Background

**GNU Taler: Introduction** 

**Protocol Basics** 

Offline payments

Future Work & Conclusion

Component Zoo

#### A Social Problem

This was a question posed to RAND researchers in 1971:

"Suppose you were an advisor to the head of the KGB. Suppose you are given the assignment of designing a system for the surveillance of all citizens and visitors within the boundaries of the USSR. The system is not to be too obtrusive or obvious. What would be your decision?"

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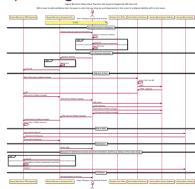
"I think one of the big things that we need to do, is we need to get away from true-name payments on the Internet. The credit card payment system is one of the worst things that happened for the user, in terms of being able to divorce their access from their identity."

—Edward Snowden, IETF 93 (2015)

### Banks have Problems, too!

3D secure ("verified by visa") is a nightmare:

- Complicated process
- Shifts liability to consumer
- Significant latency
- Can refuse valid requests
- Legal vendors excluded
- No privacy for buyers



Online credit card payments will be replaced, but with what?

#### The Bank's Problem

- ► Global tech companies push oligopolies
- Privacy and federated finance are at risk
- Economic sovereignty is in danger











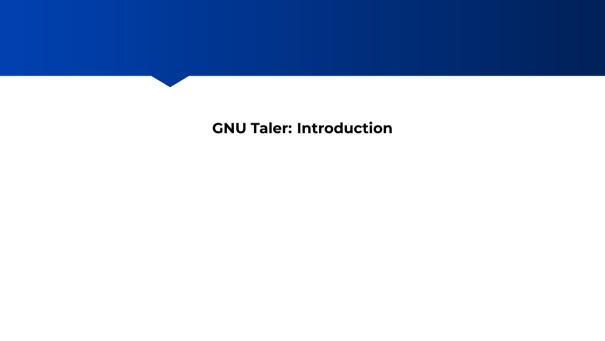
### Predicting the Future

- Google and Apple will be your bank and run your payment system
- They can target advertising based on your purchase history, location and your ability to pay
- ► They will provide more usable, faster and broadly available payment solutions; our federated banking system will be history
- After they dominate the payment sector, they will start to charge fees befitting their oligopoly size
- Competitors and vendors not aligning with their corporate "values" will be excluded by policy and go bankrupt
- ▶ The imperium will have another major tool for its financial warfare

# Central Bank Digital Currency?

Speech by Augustin Carstens, Bank of International Settlements (October 2020) on the difference between Central Bank Digital Currencies and cash.

Central Bank Digital Currency vs. Cash



# GNU Taler [1, 3, 2]

# Digital cash, made socially responsible.



Privacy-Preserving, Practical, Taxable, Free Software, Efficient

#### What is Taler?

https://taler.net/en/features.html

#### Taler is

- ▶ a Free/Libre software payment system infrastructure project
- ... with a surrounding software ecosystem
- ... and a company (Taler Systems S.A.) and community that wants to deploy it as widely as possible.

#### However, Taler is

- not a currency or speculative asset
- not a long-term store of value
- not a network or instance of a system
- not based on proof-of-work or proof-of-stake

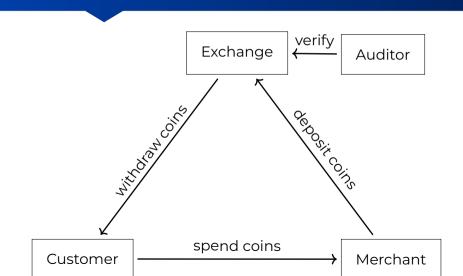
### Design principles

#### https://taler.net/en/principles.html

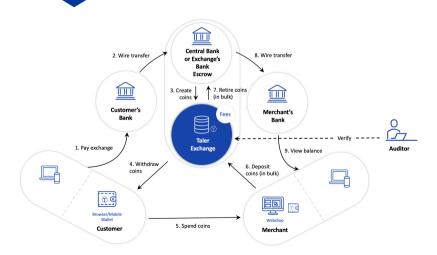
#### GNU Taler must ...

- 1. ... be implemented as free software.
- 2. ... protect the **privacy of buyers**.
- 3. ... enable the state to **tax income** and crack down on illegal business activities.
- 4. ... prevent payment fraud.
- 5. ... only disclose the minimal amount of information necessary.
- 6. ... be usable.
- 7. ... be efficient.
- 8. ... avoid single points of failure.
- 9. ... foster competition.

### **Taler Overview**



### **Architecture of Taler**



## Consumer Impact of Taler

- Convenient: pay with one click instantly in Euro, Dollar, Yen or Bitcoin
- Friction-free security: Payments do not require sign-up, login or multi-factor authentication
- Privacy-preserving: payment requires/shares no personal information
- ▶ Bank account: not required

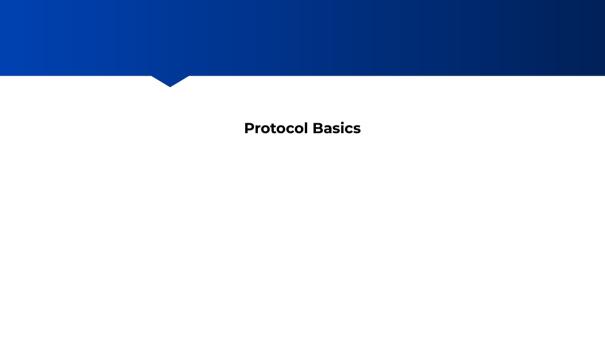
## Merchant Impact of Taler

- ► Instant clearance: one-click transactions and instant clearance at par
- ➤ Easy & compliant: GDPR & PCI-DSS compliance-free and without any effort
- Major profit increase: efficient protocol + no fraud = extremely low costs
- ▶ 1-click checkout: without Amazon and without false positives in fraud detection

# **Usability of Taler**

https://demo.taler.net/

- 1. Install browser extension.
- 2. Visit the bank.demo.taler.net to withdraw coins.
- 3. Visit the shop.demo.taler.net to spend coins.





#### How does it work?

We use a few ancient constructions:

- Cryptographic hash function (1989)
- ► Blind signature (1983)
- Schnorr signature (1989)
- ▶ Diffie-Hellman key exchange (1976) or Unique signatures (1977) or VRF (1999)
- Cut-and-choose zero-knowledge proof (1985)

But of course we use modern instantiations.

# Definition: Taxability

We say Taler is taxable because:

- Merchant's income is visible from deposits.
- Hash of contract is part of deposit data.
- State can trace income and enforce taxation.

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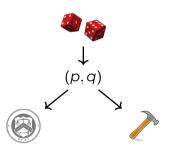
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#### Limitations:

- withdraw loophole
- sharing coins among family and friends

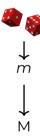
# Exchange setup: Create a denomination key (RSA)

- 1. Generate random primes p, q.
- 2. Compute n := pq,  $\phi(n) = (p-1)(q-1)$
- 3. Pick small  $e < \phi(n)$  such that  $d := e^{-1} \mod \phi(n)$  exists.
- 4. Publish public key (e, n).



# Merchant: Create a signing key (EdDSA)

- Generate random number
   m mod o as private key
- Compute public key M := mG



#### Capability:

 $m \Rightarrow$ 



# Customer: Create a planchet (EdDSA)

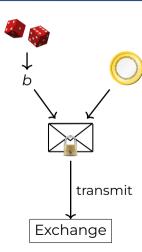
- Generate random number c mod o as private key
- ► Compute public key C := cG



Capability: c ⇒

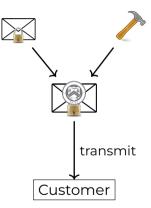
# Customer: Blind planchet (RSA)

- 1. Obtain public key (e, n)
- 2. Compute f := FDH(C), f < n.
- 3. Generate random blinding factor  $b \in \mathbb{Z}_n$
- 4. Transmit  $f' := fb^e \mod n$



# Exchange: Blind sign (RSA)

- 1. Receive f'.
- 2. Compute  $s' := f'^d \mod n$ .
- 3. Send signature s'.

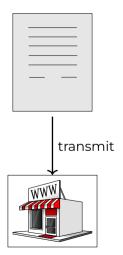


# Customer: Unblind coin (RSA)

- 1. Receive s'.
- 2. Compute  $s := s'b^{-1} \mod n$

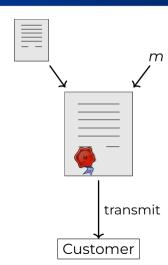


# Customer: Build shopping cart



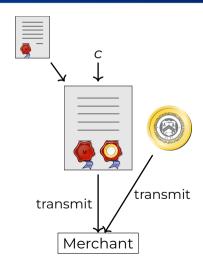
# Merchant: Propose contract (EdDSA)

- 1. Complete proposal D.
- 2. Send D,  $EdDSA_m(D)$



# Customer: Spend coin (EdDSA)

- 1. Receive proposal D,  $EdDSA_m(D)$ .
- 2. Send s, C,  $EdDSA_c(D)$



# Merchant and Exchange: Verify coin (RSA)

$$s^e \stackrel{?}{\equiv} FDH(C) \mod n$$



The exchange does not only verify the signature, but also checks that the coin was not double-spent.

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Taler is an online payment system.

# Giving change

It would be inefficient to pay EUR 100 with 1 cent coins!

- Denomination key represents value of a coin.
- Exchange may offer various denominations.
- Wallet may not have exact change!
- Must be able to pay given sufficient total funds.

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#### Method:

- ► Contract can specify to pay partial value of a coin.
- ► Allow wallet to obtain *unlinkable change*.

## Unique Signatures

- Some public key operations depend on a nonce or "random" value
  - Ex.: DSA/ECDSA (signing)
  - + same plaintext, different ciphertext
  - security may break on nonce-reuse
- Generating the nonce deterministically by hashing all inputs (see also: Fiat-Shamir transformation) can make these algorithms deterministic
  - Ex.: EdDSA
- ▶ If only one form of a valid signature exists and the verifier can check this, a signature is unique.
  - Ex.: RSA. Verifiable Random Func.

Unique signatures:





### Verifiable Random Functions

Micali, Rabin, & Vadhan (1999) proposed verifiable random functions.

Let M be some input.

- $ightharpoonup (sk, pk) := VRF_{keygen}()$
- Verifier picks M
- $ightharpoonup (v,p) := VRF_{sign}(M,sk)$
- $\triangleright$  v is deterministic, unpredictable and high-entropy for any M and sk, and (v, p) can only be computed with sk
- $\triangleright$  VRF<sub>verify</sub>(M, pk, v, p) returns true only if v was computed correctly
- sk cannot be derived from M, pk, v and p

### Straw-man solution

Given partially spent private coin key  $c_{old}$ :

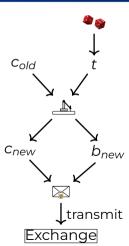
- 1. Pick random  $c_{new} \mod o$  private key
- 2. Compute  $C_{new} := c_{new}G$  public key
- 3. Pick random  $b_{new}$
- 4. Compute  $f_{new} := FDH(C_{new})$ , m < n.
- 5. Transmit  $f'_{new} := f_{new} b^e_{new} \mod n$ ... and sign request for change with  $c_{old}$ .

Cnew transmit Exchange

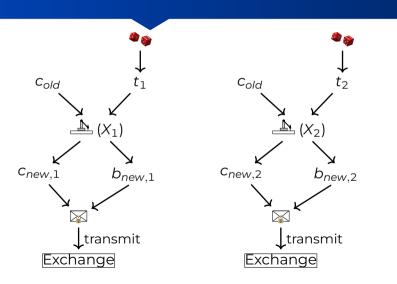
## Customer: Transfer setup (UNISIG)

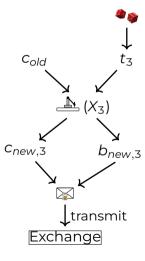
Given partially spent private coin key  $c_{old}$ :

- 1. Let  $C_{old} := c_{old}G$  (as before)
- 2. Create random nonce t
- 3. Compute unique signature  $X := UNISIG_{Cold}(t)$
- 4. Derive  $c_{new}$  and  $b_{new}$  from X using HKDF
- 5. Compute  $C_{new} := c_{new}G$
- 6. Compute  $f_{new} := FDH(C_{new})$
- 7. Transmit  $f'_{new} := f_{new} b^e_{new}$



### Cut-and-Choose





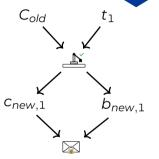
# Exchange: Choose!

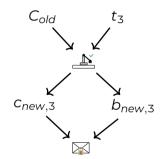
Exchange sends back random  $\gamma \in \{1, 2, 3\}$  to the customer.

### Customer: Reveal

- 1. If  $\gamma = 1$ , send  $\langle t_2, X_2 \rangle$ ,  $\langle t_3, X_3 \rangle$  to exchange
- 2. If  $\gamma = 2$ , send  $\langle t_1, X_1 \rangle$ ,  $\langle t_3, X_3 \rangle$  to exchange
- 3. If  $\gamma = 3$ , send  $\langle t_1, X_1 \rangle$ ,  $\langle t_2, X_2 \rangle$  to exchange

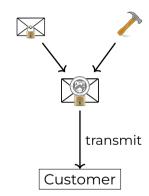
# Exchange: Verify ( $\gamma = 2$ )





# Exchange: Blind sign change (RSA)

- 1. Take  $f'_{new,\gamma}$ .
- 2. Compute  $s' := f_{new,\gamma}^{\prime d} \mod n$ .
- 3. Return signature s'.



# Customer: Unblind change (RSA)

- 1. Receive s'.
- 2. Compute  $s := s'b_{new,\gamma}^{-1} \mod n$ .

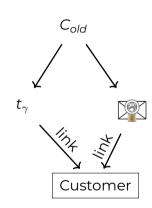


# Exchange: Allow linking change

Given Cold

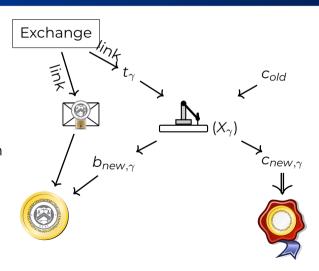
return  $t_{\gamma}$  and

 $s := s' b_{new,\gamma}^{-1} \mod n.$ 



## Customer: Link (threat!)

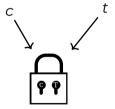
- 1. Have  $c_{old}$ .
- 2. Obtain  $T_{\gamma}$ , s from exchange
- 3. Compute  $X_{\gamma} = UNISIG_{C_{old}}(t_{\gamma})$
- 4. Derive  $c_{new,\gamma}$  and  $b_{new,\gamma}$  from  $X_{\gamma}$
- 5. Unblind  $s := s'b_{new,\gamma}^{-1} \mod n$



# VRF vs. Dold'19 with Diffie-Hellman (ECDH)

VRF/unique signatures are slightly stronger than required!

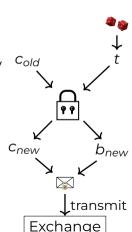
- 1. Create private keys c, t mod o
- 2. Define C = cG
- 3. Define T = tG
- 4. Compute DH cT = c(tG) = t(cG) = tC
- 5. Sign *T* with EdDSA: DH is unique, with EdDSA we have a signature, *t* allows verifier to check!



# Transfer setup with ECDH-based Refresh

Given partially spent private coin key  $c_{old}$ :

- 1. Let  $C_{old} := c_{old}G$  (as before)
- 2. Create random private transfer key t mod o
- 3. Compute T := tG
- 4. Compute  $X := c_{old}(tG) = t(c_{old}G) = tC_{old}$
- 5. Derive  $c_{new}$  and  $b_{new}$  from X
- 6. Compute  $C_{pew} := c_{pew}G$
- 7. Compute  $f_{new} := FDH(C_{new})$
- 8. Transmit  $f'_{\text{new}} := f_{\text{new}} b^{\text{e}}_{\text{new}}$



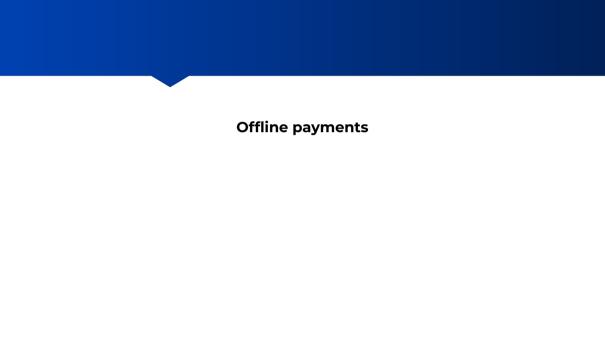
# Refresh protocol summary

- Customer asks exchange to convert old coin to new coin
- Protocol ensures new coins can be recovered from old coin
- ⇒ New coins are owned by the same entity!

#### Thus, the refresh protocol allows:

- ► To give unlinkable change.
- ► To give refunds to an anonymous customer.
- ▶ To expire old keys and migrate coins to new ones.
- To handle protocol aborts.

Transactions via refresh are equivalent to sharing a wallet.

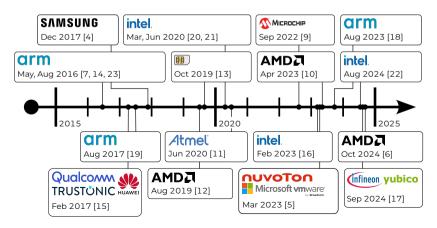


# Digitaler Euro — Offline?

Many central banks today demand offline capabilities for CBDCs.

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### A Scenario

### God is offline, but customer pays online



# Typical Payment Process All equivalent: Twint, PayPal, AliPay, PayTM

# Secure Payment ... Everything green?



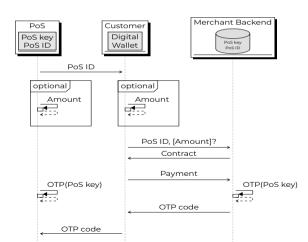
# Exploit "Code" Programming optional



# "Customers" *love* Twint ... Daily non-business for shops



# Partially Offline Payments with GNU Taler [8]

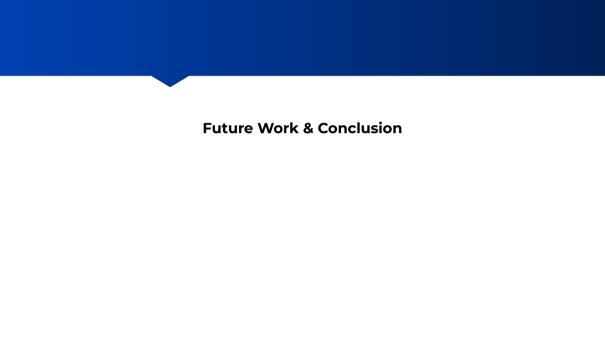


# The Emergency Act of Canada

Speech by Premier Kenney, Alberta, February 2022.

The Emergency Act of Canada

https://www.youtube.com/watch?v=NehMAj492SA(2'2022)



### Use Case: Journalism

#### Today:

- Corporate structure
- Advertising primary revenue
- Tracking readers critical for business success
- Journalism and marketing hard to distinguish

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#### With GNU Taler:

- One-click micropayments per article
- Hosting requires no expertise
- Reader-funded reporting separated from marketing
- ► Readers can remain anonymous

### Taler: Project Status

https://docs.taler.net/

- Cryptographic protocols and core exchange component are stable
- Pilot project at Bern University of Applied Sciences cafeteria
- Netzbon (regional currency) in Basel launched
- ► Taler Operations AG live Swiss-wide
- Internal alpha deployment with GLS Bank (Germany)
- Internal alpha deployment with Magnet Bank (Hungary)

# Competitor comparison

	Cash	Bitcoin	Zerocoin	Creditcard	GNU Taler
Online		++	++	+	+++
Offline	+++			+	++
Trans. cost	+			_	++
Speed	+			0	++
Taxation	_			+++	+++
Payer-anon	++	0	++		+++
Payee-anon	++	0	++		
Security	_	0	0		++
Conversion	+++			+++	+++
Libre	_	+++	+++		+++

## Other ongoing developments

- Privacy-preserving auctions (trading, currency exchange) (oezguer@taler.net)
- Hardware and software support for embedded systems (mikolai@taler.net)
- ► Tax-deductable receipts for donations to charities (donau.git)
- Unlinkable anonymous subscriptions and discount tokens (ivan@taler.net)
- ► Support for illiterate and innumerate users¹ (marc@taler.net)

<sup>&</sup>lt;sup>1</sup>Background: https://myoralvillage.org/

### **Open Challanges**

- Try to explain this to lawyers and AML staff of banks
- What are convincing arguments for citizens to switch?
- How to address anti-competitive cash-back from card payments?
- **.**..

### How to support?

```
Join: https://lists.gnu.org/mailman/listinfo/taler
Discuss: https://ich.taler.net/
Develop: https://bugs.taler.net/, https://git.taler.net/
Apply: https://nlnet.nl/propose.https://nlnet.nl/taler
Translate: https://weblate.taler.net/, translation-volunteer@taler.net
Integrate: https://docs.taler.net/
Donate: https://gnunet.org/ev
Partner: https://taler-systems.com/
```

### Conclusion

#### What can we do?

- Suffer mass-surveillance enabled by credit card oligopolies with high fees, and
- Engage in arms race with deliberately unregulatable blockchains

#### OR

Establish free software alternative balancing social goals!

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# The Taler Software Ecosystem: Overview

Taler is based on modular components that work together to provide a complete payment system:

- **Exchange:** Service provider for digital cash
  - Core exchange software (cryptography, database)
  - Air-gapped key management, real-time auditing
  - ▶ **libeufin**: Modular integration with banking systems
  - **challenger**: KYC service with OAuth 2.0 API
- ▶ **Merchant:** Integration service for existing businesses
  - Core merchant backend software (cryptography, database)
  - ▶ Back-office interface for staff
  - ► Frontend integration (E-commerce, Point-of-sale)
- ▶ Wallet: Consumer-controlled applications for e-cash
  - ► Multi-platform wallet software (for browsers & mobile phones)
  - Wallet backup storage providers (sync & Anastasis)

## Taler Exchange

The **Exchange** is the core logic of the payment system.

- One exchange at minimum must be operated per currency
- Offers a REST API for merchants and customers
- Uses several helper processes for configuration and to interact with RTGS and cryptography
- KYC support via OAuth 2.0, KycAID or Persona APIs

#### Taler Merchant

The **Merchant** is the software run by merchants to accept GNU Taler payments.

- ► REST API for integration with e-commerce
- SPA provides Web interface for administration
- ► Features include:
  - Multi-tenant support
  - Refunds
  - Templates
  - Webhooks
  - Inventory management (optional)



#### Taler Wallet

The **Wallet** is the software run by consumers to store their digital cash and authorize transactions.

- wallet-core is the logic shared by all interfaces
- Works on Android, F-Droid, iOS, Ubuntu Touch, WebExtension (Chrome, Chromium, Firefox, etc.)
- Features include:
  - Multi-currency support
  - Wallet-to-wallet payments (NFC or OR code)
  - CRDT-like data model



### **Taler Auditor**

The **Auditor** is the software run by an independent auditor to validate the operation of an Exchange.

- ► REST API for additional report inputs by merchants (optional)
- Secure database replication logic

### libeufin-nexus

libeufin-nexus allows Taler components to interact with a core banking system. It:

- provides an implementation of the Wire Gateway for the exchange
- supports EBICS 2.5 and 3.0
- other APIs such as FinTS or PSD2-style XS2A APIs can be added without requiring changes to the Exchange
- was tested with GLS Bank (DE) and Postfinance (CH) accounts and real EUR/CHF

### libeufin-bank

libeufin-bank implements a standalone bank with a Web interface. It:

- provides the Taler Core Bank API for RESTful online banking using a Web interface (with multi-factor authentication)
- includes a Taler Wire Gateway for the exchange
- offers the Taler Bank Integration API to allow wallets to easily withdraw digital cash
- optionally provides the Taler Conversion Info API for currency conversion between fiat and regional currencies
- optionally integrates with libeufin-nexus to interact with a core banking system

## Challenger

Challenger allows clients to obtain validated address (KYC) data about users:

- Customizable Web-based process for address validation
- Can validate phone numbers, e-mail addresses or physical mailing addresses
- Provides an exchange-compatible OAuth 2.0 API

## **Depolymerization**

Depolymerization is a bridge between GNU Taler and blockchains, making Taler a layer 2 system for crypto-currencies (like Lightning).

- provides an implementation of the Wire Gateway for the exchange
- Works on top of Bitcoin and Ethereum crypto-currencies, with the DLTs as the "RTGS"
- Provides same API to Exchange as libeufin-nexus

## Point-of-Sale App for Android

- Allows merchant to generate orders against Taler backend and display QR code to enable customer to pay in person
- Patterned after ViewTouch restaurant UI



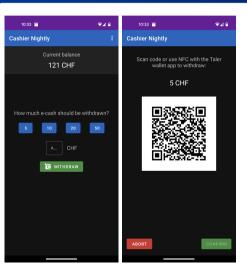
## Payment plugins



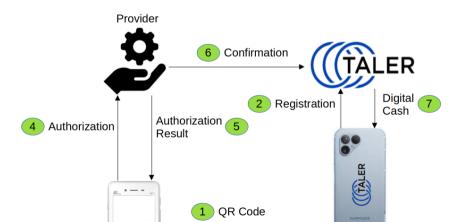
- Pretix, ticket sales system
- Joomla!, an e-commerce platform
- WooCommerce, an e-commerce solution on top of WordPress
- DrupalCommerce, an e-commerce solution on top of Drupal

## Cashier App for Android

- Enables BFH staff to convert cash to e-cash
- Staff has special bank accounts with limited funds
- Students can pay staff in cash to receive e-cash
- The Cashier App is implemented in Kotlin



## Cashless2ecash by Joel Haeberli



## TalDir (WiP)

TalDir is an extension to the existing peer-to-peer payment functionality.

- Registry to associate wallets with network addresses
- Extensible to different types of network services:
  - ► E-mail
  - ► SMS
  - Twitter
  - **.**..
- Send payments or invoices to wallets associated with network address
- ▶ Will **not** require sending wallet to use same network service

## Acknowledgments

Funded by the European Union (Project 101135475).



Co-funded by the European Union

Funded by SERI (HEU-Projekt 101135475-TALER).

#### Project funded by



Schweizerische Eidgenossenschaft Confédération suisse Confederazione Svizzera Confederaziun svizra

Swiss Confederation

Federal Department of Economic Affairs, Education and Research EAER State Secretariat for Education, Research and Innovation SERI

Views and opinions expressed are however those of the author(s) only and do not necessarily reflect those of the European Union. Neither the European Union nor the granting authority can be held responsible for them.